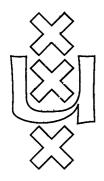
Institute for Language, Logic and Information

PHYSIOLOGICAL MODELLING USING RL

Fred de Geus Ernest Rotterdam Sieger van Denneheuvel Peter van Emde Boas

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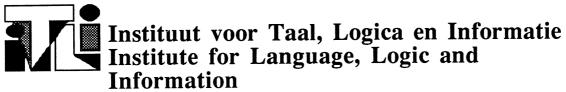
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1 Introduction

This article discusses the use of constraints to build quantitative physiological models and the application of these models to assist anaesthetists in decision making. The constraint formalism is a declarative form of knowledge representation. Each constraint embodies a model equation. A set of constraints constitutes a model. Constraint models can be extended and modified easily and the representation of model knowledge is separated from its use. The model knowledge is used by means of a constraint satisfaction mechanism. This allows the application of the same knowledge for different purposes.

In this article a simple physiological model is presented. It is represented in the relational language RL^{4,5,6,7,8} and used for two purposes: interpretation and prediction.

2 Model description

This section describes a simple physiological model of the human blood circulation and gas exchange. The model consists of 11 equations; 6 describe the systemic and pulmonary circulation and the remaining 5 the gas exchange. It is a simple model and therefore a poor reflection of reality, but it serves well to demonstrate how a model represented in constraints can be used for decision support. Figure 1 shows a graph representation of the model.

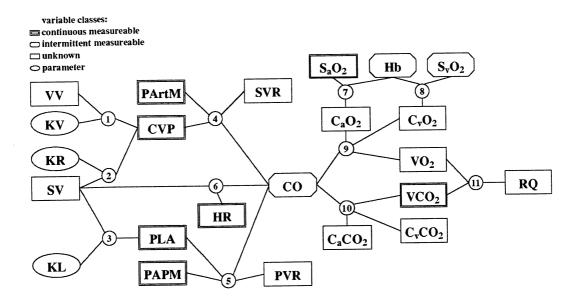


Figure 1: Graph representation of the model

Variable names used in the model are abbreviated according to convention. A summary of the properties of the variables used is shown in table 1. The numbers of the model equations listed below correspond to the numbers in the graph.

Abbre-	normal-	units^a	Name
viation	value		
CO	5500	ml/min	Cardiac Output
CVP	5-10	mmHg	Central Venous Pressure
$CaCO_2$	0.46	m ml/ml	Arterial CO ₂ Concentration
CaO_2	0.20	m ml/ml	Arterial O ₂ Concentration
$CvCO_2$	0.5	ml/ml	Venous CO ₂ Concentration
CvO_2	0.15	m ml/ml	Venous O ₂ Concentration
$_{ m HR}$	72	B/min	Heart Rate
$\mathbf{H}\mathbf{b}$	15	gr/dl	Haemoglobin concentration of blood
KL		$\mathrm{mmHg/ml}$	Left Heart parameter
KR		$\mathrm{mmHg/ml}$	Right Heart parameter
KV	_	m mmHg/ml	Venous system parameter
PAPM	13	mmHg	Mean Pulmonary Arterial Pressure
PArtM	100	\mathbf{mmHg}	Mean Arterial Pressure
PLA	6	mmHg	Left Atrium Pressure
PVR	0.0013	$\mathrm{mmHg/(ml/min)}$	Pulmonary Vascular Resistance
RQ	0.8		Respiratory Quotient
$\mathrm{SaO_2}$	100	%	Arterial O ₂ saturation
SvO_2	70	%	Venous O ₂ saturation
sv	76	ml	Stroke Volume
SVR	0.018	$\mathrm{mmHg/(ml/min)}$	Systemic Vascular Resistance
$\dot{ m VCO_2}$	200	ml/min	Expiratory CO ₂ Volume
$\dot{ m V}{ m O_2}$	250	ml/min	Inspiratory O ₂ Volume
VV	4500	ml	Venous Volume

Table 1: Properties of model variables

1. $VV = KV \times ln(CVP)$

A higher or lower venous volume (VV) results in a higher or lower central venous pressure (CVP). When VV gets larger, the veins are stretched and resist further filling. The above formula is derived from the assumption that the change in CVP due to a change in VV, is proportional to CVP: $\delta \text{CVP}/\delta \text{VV} \simeq \text{CVP}$.

2. $SV = KR \times ln(CVP)$

For the right side of the heart a similar formula can be assumed for the relationship between the stroke volume (SV) and the CVP. This formula has been chosen to keep the model simple. In reality the stroke volume increases initially, but it decreases again above a certain level of the central venous pressure. This relationship is known as the Starling curve.

3. $SV = KL \times ln(PLA)$

This relationship is the equivalent for the left side of the heart.

4. $PArtM - CVP = CO \times SVR$

Ohm's law for the systemic circulation. The pressure difference between the mean arterial pressure (PArtM) and the CVP (potential difference) is equal to the cardiac

 $^{^{}a}$ mmHg = milimeter mercury, B = beats

output (CO) multiplied by the systemic vascular resistance (SVR).

5. $PAPM - PLA = CO \times PVR$

The equivalent relationship for the pulmonary circulation.

6. $CO = SV \times HR$

The CO is the volume of one stroke multiplied by the heart rate (HR).

7. $CaO_2 = 0.000134 \times Hb \times SaO_2$

The arterial oxygen concentration (CaO_2) is proportional to the haemoglobin concentration (Hb) multiplied by the fractional arterial oxygen saturation (SaO_2) . The constant 0.000134 is in dl/g.

8. $CvO_2 = 0.000134 \times Hb \times SvO_2$

The equivalent relationship for the venous system.

9. $\dot{V}O_2 = (CaO_2 - CvO_2) \times CO$

The inspiratory oxygen volume $(\dot{V}O_2)$ is equal to the oxygen uptake, which is the difference in oxygen concentration between the arterial and the venous blood multiplied by the CO.

10. $\dot{V}CO_2 = (CvCO_2 - CaCO_2) \times CO$

The expiratory carbon dioxide volume $(\dot{V}CO_2)$ is equal to the carbon dioxide production, which is the difference in carbon dioxide concentration between the venous and arterial blood multiplied by the CO.

11. $RQ = \frac{\dot{V}CO_2}{\dot{V}O_2}$

The respiratory quotient (RQ) is defined as $\dot{V}CO_2$ divided by $\dot{V}O_2$. When no changes take place in the total body content of carbon dioxide, it has a constant value of 0.8, since the metabolic rate of oxygen used is proportional to the quantity of carbon dioxide produced.

3 Introduction to RL

The language RL^{6,7} integrates logic programming, algebraic constraint solving and relational databases. The user can combine equational constraints, Horn clauses and relational database operators to express a knowledge base. This provides maximal freedom and full conceptual transparency to the user.

Knowledge and queries expressed in RL are preprocessed by a constraint solver, and compiled into a database query. The database query is processed by an existing relational database system, so that large amounts of data can be processed efficiently.

An RL program consists of modules, each describing a system of relations. An RL query consists of a relational expression which is evaluated in the context of an existing program. RL/1 is a subset of RL. It eliminates some of the syntactic and semantic complexity, but it preserves the main idea of integrating relational databases, logic programming and algebraic constraint solving. The current implementation of RL/1 focuses on the integration of a

symbolic (numeric) constraint solver with a relational database system^{3,4,5}. A prototype system has been implemented consisting of a compiler for RL/1, a constraint solver and a relational database that is coded in Prolog. A version that produces standard SQL is currently being developed. The examples presented in the next sections were processed with this system.

In RL/1 a distinction is made between extensional and intensional objects. Extensional objects correspond to base tables in the underlying relational database. Intensional objects are defined by rules, also called clauses. A rule consists of a rule head and a rule expression separated by the keyword WHEN. The rule expression can contain system defined predicates or invocations of intensional and extensional objects, separated by a conjunctive AND operator. If an intensional object is defined by more than one rule, the rules are disjunctive. Queries in RL/1 result in an answer relation which consists of attributes equal to the attribute list between the keywords INFER and WHEN, and a (possibly empty) set of tuples. Optionally the result of a query command can be stored in a database table:

```
INFER ( attribute-list ) WHEN rule-expression [ TO table ]
```

An important feature of the RL/1 system is the capability to produce symbolic output. Processing a INFER query requires one or more invocations of the constraint solver. The solver output is then used to compile the query into a database request. The intermediate solver output can also be presented to the user directly. The syntax of a symbolic query is as follows:

```
SYMINFER ( attribute-list ) WHEN rule-expression [ TO file ]
```

The 'TO file' option allows the computed symbolic answer to be stored in a text file. Symbolic query commands result in the following answer:

```
condition= {Condition set}
solution= {Solution set}
```

Either the condition part or the solution part can be absent. The solution set contains elements of the form x = t with x a wanted variable and t a term. Wanted variables are to be eliminated from the constraint set. The condition part states under which additional restrictions the obtained solution is valid. The main goal of the constraint solver is to express the wanted variables in terms of known variables. The set of wanted variables is equal to the attribute list of the query. The known variables are determined by the rule expression of the query.

Variables occurring in the terms t of the solution set are a subset of the known variables. All wanted variables are present as one of the x variables in the solution set. The condition set consists of constraints that only contain known variables.

4 RL implementation of the model

The physiological model is represented in RL by a single rule defining the intensional object 's' as shown in Figure 2. The constraints have the same numbers as the model equations

in section 2. The '*' notation in the head of a rule indicates that all variables occurring in the associated rule expression are declared as attributes of the object. This notation can also be used in an invocation of an object and is then used to denote all attributes of the object.

In the symbolic queries below the predefined tables 'k1', 'k2', 'k3' ... will be used. If one of these tables appears in the rule expression of a query, the variables in the table become known. In addition tables are defined to group continuous measureable variables, intermittent measureable variables, non measureable variables and patient parameters. Their names are prefixed with a 'k'.

```
module phsystem(number:Ca02:CaC02:CvC02:Cv02:C0,
    number: CVP: HR: KL: KR: KV: PAPM: PArtM: PLA: PVR: RQ: Hb,
    number:Sa02:Sv02:SVR:VC02:V02:VV).
table kintermit(CO, Hb, SvO2).
table kcontinue(PArtM, CVP, PLA, PAPM, HR, Sa02, VC02).
table kunknown(VV,SV,SVR,PVR,CaO2,CaCO2,CvO2,VO2,CvCO2,RQ).
table kconstant(KV, KR, KL).
clause s(*) when
/* 1*/
         VV=KV*ln(CVP)
/* 2*/
         and SV=KR*ln(CVP)
/* 3*/
         and SV=KL*ln(PLA)
/* 4*/
         and PArtM-CVP=C0*SVR
         and APM-PLA=C0*PVR
/* 5*/
         and CO=SV*HR
/* 6*/
/* 7*/
         and Ca02=0.0001134*Hb*Sa02
/* 8*/
         and Cv02=0.0001134*Hb*Sv02
/* 9*/
         and V02=(Ca02-Cv02)*C0
         and VCO2=(CvCO2-CaCO2)*CO
/*10*/
         and RQ=VCO2/CO2.
/*11*/
close.
```

Figure 2: The physiological model in RL

5 Using the model

Representing a model in a declarative way enables its use for multiple purposes. This section discusses two ways in which a physiological constraint model can be used to assist anaesthetists in decision making: (1) the interpretation of measurements, and (2) the prediction of effects of treatment.

5.1 Interpretation of measurements

The anaesthetist tries to gain insight in the development of the physiological condition of a patient by monitoring the trend of physiological variables. Unfortunately not all variables can be measured. A physiological model, like the one presented in section 2, can be used to derive values for non measureable variables from the ones measured. The derivation of values for unmeasureable variables is a first step in the interpretation process which leads to the anaesthetists judgements and decision making.

The number of physiological variables which can be measured has increased over the past few decades. This development has raised a demand for computer systems that present all measured information in a coherent way^{1,2}. The rapid development of new measurement equipment also requires that a computer system is adaptable to new measurement possibilities.

Physiological variables can be divided into four classes. In figure 1 these categories are reflected in the shapes of variables.

- 1. continuous measureable variables: in the University Hospital of Groningen a computer system² on the theatre stores average values for these variables at 1 minute intervals. These sets of values can be used as input for interpretation.
- 2. intermittent measureable variables: values of these variables can not be automatically measured, and the measurement methods available demand extra attention of the anaesthetist. They are therefore performed only when necessary. To determine bloodgasses for example, a sample of blood has to be taken. It takes about 18 minutes for a blood sample to be analyzed in the laboratory. Another example is cardiac output, which is measured by injecting cold fluid into the blood stream. This introduces extra volume and can only be done a few times during an operation.
- 3. patient parameters: variables whose values are characteristic for a particular patient and change so slowly that they can be assumed constant with respect to other variables.
- 4. unknown variables: variables that can be neither measured nor assumed to be constant.

The possibilities for measurement interpretation offered by a physiological constraint model will be illustrated with an example: the determination of the trend of the CO. This will be followed by a more systematic approach for measurement interpretation.

5.1.1 Example: determination of the CO-trend

Anaesthetists are interested in the trend of the CO. However CO can only be measured a few times during an operation. From figure 1 it can be seen that such an incidental measurement, combined with values for the continuous measureable variables HR, CVP and PLA, determines the patient parameters KR and KL. The computation of KR and KL can be done with RL by making CO, HR, CVP and PLA known variables and KR and KL wanted. The RL query and its solution is:

KR and KL can be assumed to remain equal to the values computed for them since they are parameters. Subsequent measurements for CVP, PLA and HR can be used to compute the CO. This is done with RL by making CVP, PLA, HR, KR, and KL known variables and CO wanted. The RL query and its result is:

```
syminfer(CO) when s(*) and k5(CVP,PLA,HR,KR,KL).
condition= { KR * ln(CVP) = KL * ln(PLA) }
solution= { CO = HR * KL * ln(PLA) }
```

The solution contains an expression for the CO in terms of the known variables. The condition can be used to check whether the values measured for CVP and PLA are consistent with each other.

5.1.2 General measurement interpretation method

In the example above it was necessary to use the graph in figure 1 to decide which of the variables had to be made known and which of them wanted. The RL queries were tailored to the model and the measurements available. This is undesirable because it implies that new queries have to be developed each time the model is changed or new measurement possibilities become available. As this would practically nullify the advantages of a declarative representation, the following general method to interpret measurements with a constraint model was developed.

When measurements for a set of variables become available, they are made known to RL and all other variables are made wanted. RL produces solutions for variables that can be expressed in terms of the known variables. For the remaining non-measured variables a reduced model can be derived from the original one. In order to make RL perform the reduction, the following procedure is used:

- add equations for the available measurement values of the form 'variable_i = measurement_i' to the original model, with 'variable_i' a model variable and 'measurement_i' a value
- make the measured variables wanted and the variables for which a reduced model is searched known.

The solution derived by RL contains the added equations 'variable_i = measurement_i'. According to the specification of the RL/1 system the condition contains constraints applying to the variables declared known in the query³. These constraints make up the reduced model.

The derived relations in the reduced model can be useful for decision support. A relation between two variables for example can be displayed as an X-Y plot. When derived relations are complex however, they can not be used as such for decision support. In that case assumptions have to be made about the values of variables or about the existence of

extra constraints. Addition of assumed constraints will be demonstrated in section 5.2 on prediction of treatment effects.

It can be assumed that variables keep the value determined earlier for them. These values can be used to obtain additional variable solutions and model reductions. The four classes of variables introduced earlier can be ranked according to the validity of the assumption that their values have remained equal. This ranking is used to select the group of variables whose values are used when more solutions or further reductions are needed.

When a set of values from continuous measureable variables becomes available the sequence in which variable classes are used for value derivation and model reduction are: (1) parameter values and (2) intermittent measureable values.

For new values of intermittent measureable variables the sequence is: (1) continuous measurable values, (2) parameter values and (3) other intermittent measureable values.

Application of this general strategy results in the same values for the trend of the CO as the informal strategy that was used previously. When the value 5000 for the variable CO becomes available, RL is first requested for solutions as follows:

None of the other variables can be expressed as a function of CO, so the solution set is empty. The next step is model reduction. CO whose values is known, is made wanted to RL and the other variables are declared known:

```
syminfer(CO) when s(*) and CO=5000 and kcontinue(*) and kunknown(*)
    and kconstant(*) and k2(Hb,Sv02).
condition= {
        VC02 = RQ * V02
        CaC02 * 5000 + CvC02 * -5000 + VC02 = 0
        Ca02 * -5000 + Cv02 * 5000 + V02 = 0
        Cv02 * -7462.686567 + Hb * Sv02 = 0
        Ca02 * -7462.686567 + Hb * Sa02 = 0
        HR * SV * -0.000200 = -1
        PAPM * -0.000200 + PLA * 0.000200 + PVR = 0
        PCVP * 0.000200 + PArtM * -0.000200 + SVR = 0
        KL * ln(PLA) = SV
        ln(CVP) * KR = SV
        ln(CVP) * KV = VV
}
solution= { C0 = 5000 }
```

The reduced model is essentially the original model with 5000 substituted for the CO. It contains a relation between SV and HR which can be presented as an X-Y plot.

The next step is performed with the last set of values for the continuous measureable variables. In the next query it is assumed that these are equal to the values in table 1.

```
syminfer(VV,SV,SVR,PVR,Ca02,CaC02,Cv02,V02,CvC02,RQ,Hb,Sv02,KV,KR,KL) when
s(*) and C0=5000 and PArtM=72 and CVP=5 and PLA=17 and PAPM=27 and
HR=72 and Sa02=0.99 and VC02=200.
condition= {}
solution= { KR = 43.148259 , KL = 24.510842 }
```

When new sets of values of continuous measureable variables become available, the same three steps are performed. In the third step – the parameter substitution – the derived values for KL and KR are used. As shown below an estimate for the actual CO results. In the first query it is tried to derive values of other variables from the measurement values:

```
syminfer(VV,SV,SVR,PVR,Ca02,CaC02,Cv02,V02,CvC02,RQ,C0,Hb,Sv02,KV,KR,KL)
    when s(*) and PArtM=72 and CVP=5 and PLA=17 and PAPM=27 and
    HR=72 and Sa02=0.99 and VC02=200.
condition= {}
solution= {}
```

There are no other variables that can be derived. The solution set is empty. Next the values are used for model reduction:

```
syminfer(PArtM, CVP, PLA, PAPM, HR, SaO2, VCO2) when
    s(*) and PArtM=72 and CVP=5 and PLA=17 and PAPM=27 and
    HR=72 and Sa02=0.99 and VC02=200.
condition= {
        RQ * V02 * -0.005000 = -1
        CO * CaCO2 * O.005000 + CO * CvCO2 * -0.005000 = -1
        C0 * Ca02 * -1 + C0 * Cv02 + V02 = 0
        Cv02 * -7462.686567 + Hb * SV02 = 0
        Ca02 * -7538.067240 + Hb = 0
        CO * -0.013889 + SV = 0
        CO * PVR * -0.100000 = -1
        CO * SVR * -0.014925 = -1
        KL * -2.833213 + SV = 0
        KR * -1.609438 + SV = 0
        KV * -1.609438 + VV = 0
solution= {
        PArtM = 72
        CVP = 5
        PLA = 17
        PAPM = 27
        HR = 72
        Sa02 = 0.99
        VC02 = 200
}
```

The condition contains the reduced model. Next the parameter values are used to make further value inferences:

```
syminfer(CO,Hb,Sv02,VV,SV,SVR,PVR,Ca02,CaC02,Cv02,V02,CvC02,RQ) when
    s(*) and PArtM=72 and CVP=5 and PLA=17 and PAPM=27 and HR=72 and
    Sa02=0.99 and VC02=200 and KR=43 and KL=24.5 and KV=2750

condition= { FALSE }

solution= {
    PVR = 0.002001
    SVR = 0.013406
    SV = 69.413718
    VV = 4425.954500
    C0 = 4997.788344
}
```

The condition contains FALSE, this indicates that an inconsistency is detected. In section 5.1.1 a condition between KL, KR, CVP, PLA was derived, which is not satisfied by the given values. Inconsistencies are discussed in the next subsection.

5.1.3 Quantifying inconsistencies

RL derives a condition between the known variables as well as a solution. When the values of the measured variables do not satisfy the condition, they are inconsistent. Small inconsistencies are not alarming, since measurements always contain minor errors and models are only an abstraction of reality. The magnitudes of errors are of greater importance. They can be obtained by introducing variables for errors as follows:

- add equations of the form 'error_i = value_i variable_i' to the original model, where 'value_i' is a value which was measured or derived earlier and 'error_i' and 'variable_i' are variables,
- declare the error variables known to RL.

The RL query to find expressions relating the errors of variable values used in last query of the previous subsection is:

```
syminfer() when s(*) and EPArtM=72-PArtM and ECVP=5-CVP and EPLA=17-PLA
    and EPAPM=27-PAPM and EHR=72-HR and ESa02=.99-Sa02 and EVC02=200-VC02
    and EKL=24.5-KL and EKR=43-KR and EKV=2750-KV
    and k10(EPArtM,EVCP,EPLA,EPAPM,EHR,ESa02,EVC02,EKL,EKR,EKV).

condition= {
    ln(ECVP * -1.000000 + 5) * -1.755102 + ln(ECVP * -1.000000
    + 5) * EKR * 0.040816 + EKL * ln(EPLA * -1.000000
    + 17) * -0.040816 + ln(EPLA * -1.000000 + 17) = 0
}
```

Identifiers for error variables are formed prefixing the variable names with an E. The condition contains an expression relating the errors of KR, KL, CVP and PLA. The magnitude of the inconsistency can be obtained by substituting zero for n - 1 errors in the condition and to use RL to compute the corresponding value for the remaining error.

When the previous query is extended with:

the result is:

The error value found relative to the value of KL is: $0.073378 / 24.5 \times 100 \% = 0.3\%$.

5.2 Prediction of treatment effects

Effects of a treatment can be divided into two groups:

- primary: effects that are a direct result of the treatment
- secondary: effects due to physiological reactions provoked by the primary effects.

When an assumption is made about the primary effects of a treatment, the magnitude of its secondary effects can be computed using a physiological constraint model. The secondarily affected model variables are made the wanted variables for RL, both the primary affected variables and the ones known to be unaffected by the treatment are made known. For the primarily affected variables a new value is supplied based on knowledge about the effects of the treatment. The values supplied for the other variables are equal to their current ones. Sometimes it is necessary to add assumptions about the propagation of primary effects to secondary effects.

An example of treatment is the infusion of a plasma expander. A plasma expander has the property to remain in its entirety in the circulation; it does not migrate to the interstitium. The shift in plasma volume is nearly entirely (99.5%) in the venous compartment. The introduction of extra venous volume provokes the following chain of events:

```
\begin{array}{c} \mathrm{infusion} \to \mathrm{VV}\!\!\uparrow \to \mathrm{CVP}\!\!\uparrow \to \mathrm{SV}\!\!\uparrow \to \mathrm{CO}\!\!\uparrow \to \mathrm{PArtM}\!\!\uparrow, \mathrm{PAPM}\!\!\uparrow, \mathrm{C}_{\mathrm{V}}\mathrm{O}_{2}\!\!\uparrow, \mathrm{C}_{\mathrm{V}}\mathrm{CO}_{2}\!\!\downarrow \\ \mathrm{PArtM}\!\!\uparrow \to \mathrm{SVR}\!\!\downarrow \\ \mathrm{PAPM}\!\!\uparrow \to \mathrm{PVR}\!\!\downarrow \end{array}
```

```
(\uparrow = increases and \downarrow = decreases)
```

Physiological regulation compensates for the increase in PArtM and PAPM by a decrease of SVR and PVR. To enable computation of the new pressures and new resistances it is assumed that for the systemic circulation 60% of the change in CO is retrieved in a decrease of SVR and 40% in an increase in the pressure difference (PArtM – CVP). Lung resistance is more reactive to changes in CO; its decrease is brought about by perfusion of parts of the lung that were previously closed off. The figures for the lung circulation are 90% and 10%. These assumptions are added to the model by means of the following additional constraints:

$$m CO/CO_{old} = rate \ SVR_{old}/SVR = 60\% imes rate \ PVR_{old}/PVR = 90\% imes rate$$

 ${
m CO_{old}}$, ${
m SVR_{old}}$ and ${
m PVR_{old}}$ are fixed values, equal to the current values of the corresponding variables. Rate is a new variable. The consequences of the infusion of 1.5 liter plasma expander with all variables having their normal values as listed in table 1, can be predicted with the RL query:

```
syminfer(VO2, CVP, SV, CO, PArtM, PAPM, CvO2, CvCO2, SVR, PVR, RQ, CaO2, PLA) when
    s(*) and VV = 6000 and CO / 5500 = rate and 0.018 / SVR = 0.6 * rate
    and 0.013 / PVR = 0.1 * rate and KV = 2450 and KR = 43 and KL = 24.5
    and HR = 62 and Sa02 = 100 and CaC02 = 0.46 and Hb = 15 and Sv02 = 0.70
    and VC02 = 200.
condition= {}
solution= {
        RQ = 0.437448
        CvC02 = 0.486378
        V02 = 457.197061
        Cv02 = 0.140700
        Ca02 = 0.201000
        CO = 7582.040816
        PAPM = 145.067920
        PArtM = 176.576528
        PLA = 73.567921
        SV = 105.306122
        CVP = 11.576528
        SVR = 0.021762
        PVR = 0.009430
}
```

Some of the resulting values exceed physiological limits. This shows the limitations of the simple model used.

6 Conclusion

A considerable amount of physiological knowledge exists in algebraic form. A constraint satisfaction program enables the use of an algebraic model for various purposes.

The language RL provides a framework in which algebraic constraints can be combined with logic programming and relational databases. The prototype implementation RL/1 contains a constraint satisfaction module that generates symbolic solutions, which were used to interpret patient measurements, to predict treatment effects and to trace inconsistencies between measurements.

Substitution of measured variable values results in values for other variables and in a reduced model. When the reduced model is still too complex to provide decision support, extra assumptions have to be made. When it is assumed that variables of a particular class have kept the values that were last found for them, values of other variables can be infered.

Other forms of assumption are possible. The physician can do what-if reasoning: "if a particular value is assumed for a variable, what are the corresponding values for other

variables?" The physician can also add extra constraints to the model. This was shown in the prediction of treatment effects.

Interaction with the anaesthetist is required about intermediate results and further assumptions to be made. As time in the operating theatre is limited, an optimal interface is essential. Effective interaction must be possible however, because making assumptions is the way in which physicians cope with incompleteness of information. A constraint solver provides a tool to support this.

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