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PROVABLE FIXED POINTS IN $I\Delta_0+\Omega_1$

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1. INTRODUCTION [1]

This work should be considered as part of the general investigation into the arithmetical system $I\Delta_0+\Omega_1$. We will present a refinement to $I\Delta_0+\Omega_1$ of a result given in [deJongh-Montagna, 1988], on witness comparison formulas having only provable fixed points in PA.

Briefly, let us introduce the arithmetical system and some of its properties: $I\Delta_0+\Omega_1$ (Cf. [Paris-Wilkie, 1987]) is a set of axioms expressing the elementary arithmetic properties of the basic symbols 0, ', +, *, ≤ (in the following we will refer to the obvious first order language containing these symbols as $\mathfrak{S}$) together with the bounded induction schema $I\Delta_0$ (defined in $\mathfrak{S}$):

$$\forall x,z\ (\varphi(x,0) \wedge \forall y{\le}z.\ (\varphi(x,y) \rightarrow \varphi(x,y')) \rightarrow \forall y{\le}z\ \varphi(x,z)) \qquad (\varphi\in\Delta_0)$$

plus an $\mathfrak{S}$-sentence expressing $\forall x\exists y\ \omega_1(x)=y$ where $\omega_1(x):= x^{|x|}$ and $|-|$ is the length function for the binary representation of x.

We note that by the following result of [Verbrugge,1989]

If NP≠CO-NP then

$$I\Delta_0+\Omega_1 \nvdash \forall b,c(\exists a\ (Prf(a,c) \wedge \forall z{\le}a\ \neg Prf(z,b)) \rightarrow Pr(\ulcorner \exists a Prf(a,c) \wedge \forall z{\le}a \neg Prf(z,b)\urcorner))$$

it seems highly unlikely that the principle of *Σ_1-completeness* , i.e.

$$\varphi \rightarrow Pr(\ulcorner\varphi\urcorner) \text{ for } \varphi\in\Sigma_1$$

is provable in $I\Delta_0+\Omega_1$. However, it can be shown that $I\Delta_0+\Omega_1$ proves *Švejdar's principle* (Cf.[Švejdar,1983]): i.e.

$$I\Delta_0+\Omega_1 \vdash Pr(\ulcorner\varphi\urcorner)\rightarrow Pr(\ulcorner\exists a(Prf(a,\ulcorner\psi\urcorner)\wedge\forall z{\le}a\neg Prf(z,\ulcorner\varphi\urcorner)) \rightarrow \psi\urcorner) \text{ (for all } \varphi,\psi)$$

(Cf. [Verbrugge,1989]) and

$$I\Delta_0+\Omega_1 \vdash Pr(\ulcorner C(S) \rightarrow s'\urcorner) \rightarrow Pr(\ulcorner s'\urcorner)$$

where $C(S)= \bigwedge\{ s \rightarrow Pr(\ulcorner s\urcorner) : s\in S.\}$, S is a finite set of

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[1] Prerequisites: the reader is supposed to be familiar with [Smoryński,1985]; knowledge of [de Jongh-Montagna,1988] will be helpful.

Σ_1-sentences and s' is a Σ_1-sentence (Cf. [Visser,1989]). In the following this last principle will be called *Visser's principle*.

In [Paris-Wilkie,1987], [Buss,1986] and [Verbrugge,1989] ample motivation for the general study of $I\Delta_0+\Omega_1$ is given; therefore we will turn our attention directly to the more specific aim of this paper.

In [Parikh,1971] it is shown that for each primitive recursive function g, there is a Σ_1-formula s such that PA $\vdash$ s and

$$g(\mu z.Prf_{PA}(z,\ulcorner Pr_{PA}(\ulcorner s\urcorner)\urcorner)) < \mu z.Prf_{PA}(z,\ulcorner s\urcorner) \qquad (*)$$

The result is based on the fact that (*) has *only provable fixed points.*

In [deJongh-Montagna, 1988] Parikh's result is analyzed in the modal context R (Cf. [Guaspari-Solovay, 1979]) when g is the identity function; a characterization is given for pairs of modal formulas B(p) and C(p) such that for each arithmetical interpretation *,

if PA $\vdash p^* \leftrightarrow (\Box B(p) \prec \Box C(p))^*$ then PA $\vdash p^*$: $\Box B(p) \prec \Box C(p)$ has only provable fixed points in PA. In [deJongh-Montagna, 1989] the result is extended to arbitrary g which are provably recursive in PA.

Our aim is to refine the positive part of the proof of [deJongh-Montagna, 1988], the part in which it is shown that the formulas specified do indeed have only provable fixed points in PA, to a weaker modal system in which the *Σ-completeness axiom* (i.e. the corresponding modal version of the Σ_1-completeness principle) does not hold.
In section 3, it is shown that the modal version of Visser's principle: i.e.

$$(V) \qquad \Box(C(S) \rightarrow s') \rightarrow \Box s'$$

where $C(S) = \bigwedge\{ s \rightarrow \Box s : s\in S\}$,
S is a finite set of Σ-formulas
s' is a Σ-formula,

playing the role of a weak version of Σ-completeness, suffices to obtain the refined theorem we are looking for.
What is provable in the weak modal system including Visser's schema, is clearly provable in $I\Delta_0+\Omega_1$ under every arithmetical interpretation; therefore, it follows that PA has no witness comparison formulas having only provable fixed points which the system $I\Delta_0+\Omega_1$ does not already have.

Based on the result obtained in section 3, in section 4 we present a counterexample to show that the modal version of Švejdar's principle

$$(\check{S}v) \qquad \Box A \to \Box(\Box B \preccurlyeq \Box A \to B) \qquad \text{for all formulas A,B}$$

does not imply Visser's schema: the result will give an insight to understand why Švejdar's schema cannot play much of a role in the study of formulas having only provable fixed points.
In an appendix we give some proofs, mainly due to Visser [1989], of modal principles derivable from Visser's principle.

2. MODAL SYSTEMS AND KRIPKE SEMANTICS

In this section we will briefly introduce the modal systems that we are going to work with, together with the associated Kripke-semantics.

Formulas of our system are built up from propositional atoms using the boolean connectives $\wedge, \vee, \neg, \to, \leftrightarrow, \top, \bot$, a unary modality $\Box$ and binary witness comparisons $\prec$, $\preccurlyeq$, where $\prec$ and $\preccurlyeq$ are applicable only to those formulas having $\Box$ as principal connective. The following definition will introduce the list of modal systems.

Def 2.1:

(a) B^- (*Basic System*) is the modal system L (Prl in [Smoryński, 1985]) (including its rules: modus ponens and necessitation) to which the following *order axioms* are added (Cf. [deJongh, 1987]):

$$\text{(O1)} \quad \Box A \to (\Box A \preccurlyeq \Box B \vee \Box B \prec \Box A)$$
$$\text{(O2)} \quad \Box A \preccurlyeq \Box B \to \Box A$$
$$\text{(O3)} \quad \Box A \preccurlyeq \Box B \wedge \Box B \preccurlyeq \Box C \to \Box A \preccurlyeq \Box C$$
$$\text{(O4)} \quad \Box A \prec \Box B \leftrightarrow (\Box A \preccurlyeq \Box B \wedge \neg(\Box B \preccurlyeq \Box A))$$

(b) Z^- (Cf. [Švejdar, 1983]) is the system B^- plus Švejdar's schema:

$$(\check{S}v) \qquad \Box A \to \Box(\Box B \preccurlyeq \Box A \to B) \qquad \text{for all formulas A,B}$$

(c) BV^- is the system B^- plus Visser's schema:

$$(V) \qquad \Box(C(S) \to s') \to \Box s'$$

where $C(S) = \bigwedge\{ s \to \Box s : s \in S\}$, S is a finite set of Σ-formulas and s' is a Σ-formula

(d) B, BV, Z are respectively the systems B^-, BV^- and Z^- with the rule $\Box$E:

$$\Box A/A \qquad \text{(for all formulas A)}$$

added.

Let A(p) be some formula of B of the form $\Box B(p) \preccurlyeq \Box C(p)$. As in [deJongh-Montagna,1988] we take BC^-, BVC^- and ZC^- to be the systems B^-, BV^- and Z^- respectively, plus the axiom $c \leftrightarrow A(c)$ (analogous notation is used for the systems B, BV and Z). Since a different system is defined for different choice of A it would be more appropriate to name the systems $BC(A)^-$, $BVC(A)^-$ and $ZC(A)^-$. But, as it will always be clear in the sequel which formula A is intended, we will refrain from doing so, in order not to unnecessarily complicate the notation.

Def 2.2: a *model for* B^- is a finite, tree-ordered Kripke-model for L in which witness comparison formulas are treated as atomic formulas and in which every instance of (O1)-(O4) is forced at each node.

Def 2.3: *models for* BV^-, Z^- are Kripke-models for B^- where respectively (V), (Šv) is forced at each node.

It is appropriate to remark that, just as is pointed out in [Verbrugge,1989] for the system Z^-, also for BV^- the forcing for witness comparison formulas in BV^--Kripke-models is *not persistent*, i.e. it does not necessarily hold that if $j \Vdash \Box A \preccurlyeq \Box B$ (resp. $j \Vdash \Box A \prec \Box B$) and jRk then $k \Vdash \Box A \preccurlyeq \Box B$ (resp. $k \Vdash \Box A \prec \Box B$).

Def 2.4: a Kripke-model is *A-sound* if its root satisfies $\Box B \to B$ for every subformula B of A.

Theorem 2.5: let A be a modal formula; then
(i) $\vdash_{BV} A \Rightarrow$ A is valid in every A-sound Kripke model for BV^-
(ii) $\vdash_{Z} A \Leftrightarrow$ A is valid in every A-sound Kripke model for Z^-

Pf: (i) left to the reader; (ii) in [Švejdar,1983]. ♦

Corollary 2.6: let A be a modal formula.

(i) $\vdash_{BVC} A \Rightarrow$ A is valid in every A-sound Kripke model for BV^- where $c \leftrightarrow A(c)$ is forced at every node.

(ii) $\vdash_{ZC} A \Rightarrow$ A is valid in every A-sound Kripke model for Z^- where $c \leftrightarrow A(c)$ is forced at every node.

Pf: left to the reader. ♦

3. WITNESS COMPARISON FORMULAS HAVING ONLY PROVABLE FIXED POINTS IN BV

Theorem 3.3 of [deJongh-Montagna,1988] reads:

> If B(p) and C(p) are L-formulas (i.e. do not contain witness comparisons), possibly containing propositional variables other than p, then $A(p) \equiv \Box B(p) \preccurlyeq \Box C(p)$ has only provable fixed points in R iff
> (i) $\vdash_L B(T)$
> (ii) $\vdash_L \Box^+(\Box B(\bot) \to \Box C(\bot)) \to \Box^{k+1}\bot$, for some k
> ($\Box^+ D$ abbreviates $D \wedge \Box D$)

Our aim is to obtain a characterization for a witness comparison formula to have only provable fixed points in BV. The result presented in this section constitutes a refinement of the theorem proved by de Jongh and Montagna; the proof that we present is syntactical and based on a different approach characterized by the proof of the following theorem:

Theorem 3.1: Let B(p) and C(p) be L-formulas. If
(i) $\vdash_L B(T)$
(ii) $\vdash_L \Box^+(\Box B(\bot) \to \Box C(\bot)) \to \Box^{k+1}\bot$, for some k,
then $A(p) \equiv \Box B(p) \preccurlyeq \Box C(p)$ has only provable fixed points in BV.

Some preparatory lemmas are needed. In the following we assume that (i) and (ii) of theorem 3.1 hold, the systems BC^-, BVC^- and BVC refer to the A(p) of this theorem. Some results already proved by Visser (Cf. [Visser,1989]) for his principle and used in the proof of the following lemmas are given in the appendix.

Lemma 3.2: $\vdash_{BC^-} \Box^+\neg c \to \Box^{k+1}\bot$

Pf:

1. $\vdash_{B^-} \Box\neg c \to \Box(c \leftrightarrow \bot)$
$\to \Box^+((\Box B(c) \leftrightarrow \Box B(\bot)) \wedge (\Box C(c) \leftrightarrow \Box C(\bot)))$
$\to (\Box^+(\Box B(c) \to \Box C(c)) \to \Box^{k+1}\bot)$ (by (b) and the Substitution Lemma (Cf. [Smoryński, 1985]))
2. $\vdash_{BC^-} \Box^+\neg c \to \Box^+(\Box B(c) \to \Box C(c))$ (by obvious properties of $\preccurlyeq$)
3. $\vdash_{BC^-} \Box^+\neg c \to \Box^{k+1}\bot$ (by 1 and 2). ♦

Lemma 3.3: $\vdash_L \Box c \to \Box B(c)$

Pf:

1. $\vdash_L c \to B(T)$ (by (i))
2. $\vdash_L \Box c \to \Box B(T)$
3. $\vdash_L \Box c \to \Box(c \leftrightarrow T)$
4. $\vdash_L \Box c \to \Box B(c)$ (by 2 and 3). ♦

Lemma 3.4: $\vdash_L \Box^+ c \to \Box^+ B(c)$

Pf:

1. $\vdash_L c \to B(T)$ (by (i))
2. $\vdash_L \Box^+ c \to \Box^+ B(T)$
3. $\vdash_L \Box^+ c \to \Box^+(c \leftrightarrow T)$
$\to (\Box^+ B(c) \leftrightarrow \Box^+ B(T))$
4. $\vdash_L \Box^+ c \to \Box^+ B(c)$ (by 2 and 3). ♦

Lemma 3.5: $\vdash_L \Box^{k+1}\bot \to (\Box C(\bot) \to B(\bot))$

Pf: We claim that, if $\vdash_L \Box^+(\Box B \to \Box C) \to \Box^{k+1}\bot$, then $\vdash_L \Box^{k+1}\bot \to (\Box C \to B)$, where B,C are arbitrary L-formulas.
For suppose not, then a model **M** exists such that
$\mathbf{M} \models \Box^+(\Box B \to \Box C) \to \Box^{k+1}\bot$ and $w \Vdash \Box^{k+1}\bot \wedge \Box C$, $w \nVdash B$, for some node w in **M**. Take the submodel of **M** generated by w and add a tail of nodes below w of such a length that the new model gets a root x of level greater than or equal to k+1 (end nodes are counted as having level 0). Clearly none of the nodes added below w can force $\Box B$ but all of them force $\Box^+(\Box B \to \Box C)$. By hypothesis, $x \Vdash \Box^{k+1}\bot$ and this gives a contradiction, which proves our claim.

By the claim and (ii) it follows that: $\vdash_L \Box^{k+1}\bot \to (\Box C(\bot) \to B(\bot))$. ♦

Lemma 3.6: $\vdash_{BC^-} \Box^+\neg c \rightarrow \Box^+B(c)$

Pf:

1. $\vdash_L \Box^+\neg c \rightarrow (\Box C(c) \leftrightarrow \Box C(\bot)) \wedge (B(c) \leftrightarrow B(\bot))$
2. $\vdash_{BC^-} \Box^+\neg c \rightarrow (\Box C(c) \rightarrow B(c))$ (by lemma 3.2 and lemma 3.5)
3. $\vdash_{BC^-} \neg c \rightarrow (\Box B(c) \rightarrow \Box C(c))$ (by obvious properties of $\preccurlyeq$)
4. $\vdash_{BC^-} \Box^+\neg c \rightarrow (\Box B(c) \rightarrow B(c))$ (by 2 and 3)
5. $\vdash_{BC^-} \Box^+\neg c \rightarrow \Box(\Box B(c) \rightarrow B(c))$
   $\rightarrow \Box B(c)$ (by formalized Löb)
6. $\rightarrow B(c)$ (by 4) ♦

Lemma 3.7: $\vdash_{BV^-} \Box\Box A \vee \Box\Box B \rightarrow \Box(\Box^+(\Box A \prec \Box B) \vee \Box^+(\Box B \preccurlyeq \Box A))$

Pf:

1. $\Box\Box A \vee \Box\Box B \rightarrow \Box(\Box A \prec \Box B \vee \Box B \preccurlyeq \Box A)$
   $\rightarrow \Box((\Box A \prec \Box B \rightarrow \Box(\Box A \prec \Box B) \wedge \Box B \preccurlyeq \Box A \rightarrow \Box(\Box B \preccurlyeq \Box A)) \rightarrow$
   $\rightarrow (\Box^+(\Box A \prec \Box B) \vee \Box^+(\Box B \preccurlyeq \Box A)))$
   $\rightarrow \Box(\Box^+(\Box A \prec \Box B) \vee \Box^+(\Box B \preccurlyeq \Box A))$ (by (V)) ♦

Corollary 3.8: $\vdash_{BV^-} \Box A \vee \Box B \rightarrow \Box(\Box A \prec \Box B \rightarrow \Box(\Box A \prec \Box B))$

Pf: Trivial. ♦

Lemma 3.9: $\vdash_{BVC^-} \Box^{k+2}c \rightarrow \Box^n\Box B(c)$ for each $0 \leq n \leq k+1$

Pf: by downward induction on n:

$n=k+1$: $\vdash_L \Box^{k+2}c \rightarrow \Box^{k+1}\Box c$
$\rightarrow \Box^{k+1}\Box B(c)$ by lemma 3.3;

$n<k+1$: recall that by induction hypothesis we have
$\vdash_{BVC^-} \Box^{k+2}c \rightarrow \Box^{n+1}\Box B(c)$, i.e.
$\vdash_{BVC^-} \Box^{k+2}c \rightarrow \Box^n\Box\Box B(c)$, So,

1. $\vdash_{BV^-} \Box^{k+2}c \rightarrow (\Box^{n+1}\Box B(c) \rightarrow$
   $\rightarrow \Box^{n+1}(\Box^+(\Box B(c) \preccurlyeq \Box C(c)) \vee \Box^+(\Box C(c) \prec \Box B(c)))$
   (by lemma 3.7)
2. $\vdash_{BVC^-} \Box^{k+2}c \rightarrow (\Box^{n+1}\Box B(c) \rightarrow \Box^{n+1}(\Box^+c \vee \Box^+\neg c))$
   $\rightarrow (\Box^{n+1}\Box B(c) \rightarrow \Box^{n+1}B(c))$ (by lemma 3.4 and lemma 3.6)
   $\rightarrow \Box^{n+1}B(c)$ (by modus ponens with the induction hypothesis). ♦

We are now ready to prove theorem 3.1:

Pf (theorem 3.1):

1. $\vdash_{BVC^-} \Box^{k+2}c \to \Box B(c)$ (by lemma 3.9 where n=0)

   $\to \Box B(c) \preccurlyeq \Box C(c) \vee \Box C(c) \prec \Box B(c)$ (by obvious properties of $\preccurlyeq$)

2. $\vdash_{BVC^-} \Box(\Box B(c) \preccurlyeq \Box C(c)) \vee \Box(\Box C(c) \prec \Box B(c)) \to \Box^+ c \vee \Box^+ \neg c$

   $\to \Box c \vee \Box^+ \neg c$

   $\to \Box^{k+1} c \vee \Box^{k+1} \bot$

   (by lemma 3.2)

   $\to \Box^{k+1} c$

3. $\vdash_{BVC^-} \Box^{k+1} c$ (by 1, 2)

4. $\vdash_{BVC}$ c (by $\Box$E) ♦

The refinement that we were looking for is an immediate consequence of theorem 3.1:

Theorem 3.10: Let B(p) and C(p) be L-formulas; then
$A(p) \equiv \Box B(p) \preccurlyeq \Box C(p)$ has only provable fixed points in BV
iff
(i) $\vdash_L B(\top)$
(ii) $\vdash_L \Box^+(\Box B(\bot) \to \Box C(\bot)) \to \Box^{k+1}\bot$, for some k.

Pf: ($\Rightarrow$) If c is a fixed point for A(p) then $\vdash_{BVC}$ c, therefore $\vdash_{RC}$ c and by lemma 2.3 in [deJongh-Montagna, 1988]
$\vdash_{R^-} \Box^+(c \leftrightarrow A(c)) \to \Box^{k+1} c$ for some k. Now apply theorem 3.3 in [deJongh-Montagna,1988].
($\Leftarrow$) by theorem 3.1. ♦

By theorem 3.10 and theorem 3.3 (Cf. [deJongh-Montagna,1988]) it follows that the formulas of the form $A(p) \equiv \Box B(p) \preccurlyeq \Box C(p)$ having only provable fixed points in R are exactly the formulas having only provable fixed points in BV. In other words, to obtain the formulas having only provable fixed points we do not need the strong Σ-completeness schema (i.e. $A \to \Box A$, for every Σ-formula A) but we can replace it by the weaker (V).

Although theorem 3.10 is formulated with *iff* one should note that, unlike with R and PA, $A(p) \equiv \Box B(p) \preccurlyeq \Box C(p)$ having only provable fixed points in $I\Delta_0+\Omega_1$ for all arithmetical interpretations does not imply that A(p) has only provable fixed points in BV, since arithmetic completeness even of L is unknown for $I\Delta_0+\Omega_1$ (see [Verbrugge,1989]). At the present, theorem 3.9 does imply that each formula of R having only provable fixed points in PA has only provable fixed points in $I\Delta_0+\Omega_1$ when arithmetical interpretations are restricted to sentences. The restriction to sentences is essential; otherwise Visser's principle loses its validity (see [Visser, 1989]).

4. INDEPENDENCE OF VISSER'S AND ŠVEJDAR'S SCHEMAS

In this section we will prove that Švejdar's schema does not imply Visser's schema. To show that, consider the formula $\Box^3 p \prec \Box^2 p$ having only provable fixed points in R, as proved in [deJongh-Montagna,1988]. By theorem 3.9 it follows that this formula has only provable fixed points in BV. On the other hand[2], notice that $\Box^3 p \prec \Box^2 p$ cannot have only provable fixed points in Z because by Švejdar's faithful interpretation of $\Box A \prec \Box B$ as "there exists a proof of A using axioms with smaller Gödel numbers than any proof of B" (Cf. [Švejdar,1983]) that would mean that for the fixed point c in PA, $\Box^2 c$ would have a proof in PA using less axioms than any proof of $\Box c$ would use. This is impossible because being a provable Σ-sentence, $\Box c$ wouldn't need any but the axioms of Q and we could take those as the zero base of our interpretation. This proves our claim.

At this point it may be of interest to remark that the formula $\Box^2 p \prec \Box p$ has only provable fixed points in Z. The following argument is due to Visser: In BC^- it is provable that $\Box^2 c \to \Box(\Box c \preccurlyeq \Box^2 c \vee \Box^2 c \prec \Box c)$. Thus, in ZC^-, $\Box^2 c \to \Box c$ is provable, from which with Löb, immediately c follows. Under the same arithmetical interpretation used in the previous argument, the result is not very surprising: it is well known that there are theorems provable in PA and not in Q. From these observation we can see that Švejdar's schema can hardly be useful in studying formulas having only provable fixed points. Recall also that in the proof of theorem 3.9, the schema (Šv) is not used.

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[2] the argument was suggested to the author by F.Montagna.

APPENDIX: SOME THEOREMS PROVED BY (V)

In [Visser,1989] the following theorems, proved using the principle (V), are pointed out:

(V1) $\Box \bigvee S \rightarrow \Box \bigvee S^+$

(V2) $\Box(\Box A \rightarrow \bigvee S) \wedge \Box(\bigvee S^+ \rightarrow A) \rightarrow \Box A$

(V3) $\Box(C(S) \rightarrow (A \rightarrow s')) \rightarrow \Box A \rightarrow \Box s'$

(V4) $\Box(C(S) \rightarrow (\Box s' \rightarrow s')) \rightarrow \Box s'$

where S is a finite set of Σ-formulas, $C(S) = \bigwedge\{ s \rightarrow \Box s : s \in S\}$, $S^+ = \{ s \wedge \Box s : s \in S\}$ and s' a Σ-formula .

We will give the proof of them in the modal system BV^-:

(V1):

1. $\Box \bigvee S \rightarrow \Box(C(S) \rightarrow \bigvee \Box^+ S)$
2. $\Box(C(S) \rightarrow \bigvee \Box^+ S) \rightarrow \Box(\bigvee \Box^+ S)$ (by (V))
3. $\Box \bigvee S \rightarrow \Box(\bigvee \Box^+ S)$ (by 1 and 2)

(V2):

1. $\Box(\Box A \rightarrow \bigvee S) \rightarrow \Box(\Box\Box A \rightarrow \Box \bigvee S)$

   $\rightarrow \Box(\Box\Box A \rightarrow \Box(\bigvee \Box^+ S))$ (by (V1))
2. $\Box(\bigvee \Box^+ S \rightarrow A) \rightarrow \Box(\Box \bigvee \Box^+ S \rightarrow \Box A)$
3. $\Box(\Box A \rightarrow \bigvee S) \wedge \Box(\bigvee \Box^+ S \rightarrow A) \rightarrow \Box(\Box\Box A \rightarrow \Box A)$ (by 1 and 2)

   $\rightarrow \Box\Box A$ (by formalized Löb)

   $\rightarrow \Box \quad S$

   $\rightarrow \Box(\bigvee \Box^+ S)$ (by (V1))

   $\rightarrow \Box A$

(V3):

1. $\Box(C(S) \rightarrow (A \rightarrow s')) \rightarrow \Box(A \rightarrow (C(S) \rightarrow s'))$

   $\rightarrow \Box A \rightarrow \Box(C(S) \rightarrow s')$

   $\rightarrow \Box A \rightarrow \Box s'$ (by (V))

(V4):

1. $\Box(C(S) \rightarrow (\Box s' \rightarrow s')) \rightarrow \Box(\Box(C(S) \rightarrow (\Box s' \rightarrow s')))$

$\rightarrow \Box(\Box\Box s' \rightarrow \Box s')$ (by (V3))

$\rightarrow \Box\Box s'$ (by formalized Löb)

$\rightarrow \Box(C(S) \rightarrow \Box s')$

$\rightarrow \Box(C(S) \rightarrow s')$

$\rightarrow \Box s'$ (by (V))

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